

Theory of Noiseless Subsystems

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Outline

Noiseless Subsystems (NSs)

• *What they are –*

Most general route to the emergence of **noise-protected** degrees of freedom in open quantum systems.

Knill, Laflamme, Viola, PRL **84**, 2525 (2000).

• *What they are good for –*

Most general structure for noise-protected **quantum information** storage. Strategies to quantum noise control are understandable in terms of identification/generation of appropriate NSs.

• *Where do we stand in practice –*

- **Quantum optics**, 2 qubits: Preservation of a 1D subspace (**singlet**) against collective noise [Kwiat et al., *Science* **290**, 498 (2000)]
- **Trapped ions**, 2 qubits: Preservation of a 2D subspace against collective **dephasing** [Kielpinski et al., *Science* **291**, 1013 (2001)]
- **NMR**, 3 qubits: Preservation of a 2D subsystem against **general** collective noise [In progress, MIT Spatial NMR Laboratory]

Noiseless sub-spaces vs sub-systems

Interaction between system S and environment E

$$H_{SE} = \sum_{\alpha} E_{\alpha}^{(S)} \otimes B_{\alpha}^{(E)}$$

Goal: Noiseless one-qubit coding

- **Noiseless subspaces** – Look for a degenerate action of the interaction operators on a subspace of states of S:

$$E_{\alpha}^{(S)} |i_L\rangle = c_{\alpha} |i_L\rangle, \quad \forall \alpha, i=0,1$$

States $|i_L\rangle\langle i_L|$ are pure, preserved states of S.

- **Noiseless subsystems** – Look for a degenerate action of the interaction operators on a given quantum number (“factor”) of a subspace of states of S:

$$E_{\alpha}^{(S)} (|\lambda_L\rangle \otimes |m_Z\rangle) = c_{\alpha} |\lambda_L\rangle \otimes \Omega_{\alpha}^{(S)} |m_Z\rangle, \quad \forall \alpha, \lambda=0,1$$

States $|\lambda_L\rangle\langle \lambda_L| \otimes \rho^{(Z)}$ need not be pure states of S.

Quantum **information** stored in the L-subsystem is preserved.

From three symmetric spins ...

System S: Three distinguishable spin 1/2 particles

$$S = \text{span} \{ \{ |0\rangle_1, |1\rangle_1 \} \otimes \{ |0\rangle_2, |1\rangle_2 \} \otimes \{ |0\rangle_3, |1\rangle_3 \} \} \simeq \mathbb{C}^8$$

Environment E couples **symmetrically** to each spin:

$$E_\alpha^{(S)} = \left(\sigma_\alpha^{(1)} + \sigma_\alpha^{(2)} + \sigma_\alpha^{(3)} \right) / 2 = J_\alpha, \quad \alpha = x, y, z$$

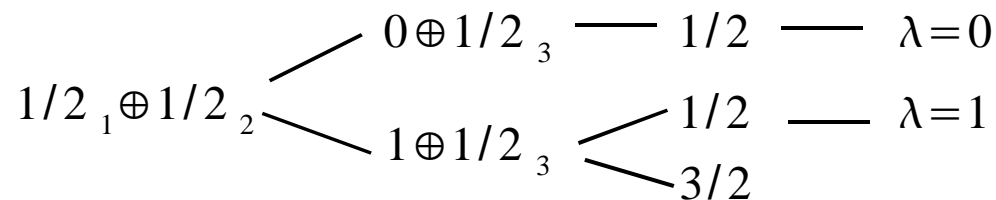
- There are **no** noiseless subspaces of S.

Can we identify a protected degree of freedom in S ?

- Choose on S a basis of joint J^2, J_z -eigenstates:

$$S = S_{3/2} \oplus S_{1/2}, \quad \dim S_J = 4$$

$$S_{1/2} = \text{span} \{ |\lambda, j_z\rangle_{1/2} \mid \lambda = 0, 1; j_z = \pm 1/2 \}$$



Keyword: The J_α have an **identity action** on the coordinate λ .

... to a noiseless qubit

Write $S_{1/2} \simeq S_L \otimes S_Z$ via $|\lambda, j_z\rangle_{1/2} \simeq |\lambda\rangle_L \otimes |j_z\rangle_Z$

- The action of the noise operators on $S_{1/2}$ takes the form

$$J_\alpha \simeq \mathbf{1}^{(L)} \otimes \Omega_\alpha^{(Z)}, \quad \Omega_\alpha^{(Z)} \in \text{Mat}(2 \times 2, \mathbb{C})$$

Ω_α depends on the choice of basis states $\{|\lambda\rangle_L\}$ in S_L .

- An explicit realization:

$$|0\rangle_L \otimes |+1/2\rangle_Z = \frac{1}{\sqrt{3}} (|001\rangle + \omega |010\rangle + \omega^2 |100\rangle)$$

$$|1\rangle_L \otimes |+1/2\rangle_Z = \frac{1}{\sqrt{3}} (|001\rangle + \omega^2 |010\rangle + \omega |100\rangle)$$

$$|0\rangle_L \otimes |-1/2\rangle_Z = \frac{1}{\sqrt{3}} (|110\rangle + \omega |101\rangle + \omega^2 |011\rangle)$$

$$|1\rangle_L \otimes |-1/2\rangle_Z = \frac{1}{\sqrt{3}} (|110\rangle + \omega^2 |101\rangle + \omega |011\rangle)$$

$$\omega = e^{2\pi i/3}$$

- The factor S_L supports a **noiseless subsystem** of S .

S_L is the state space of a **noiseless qubit**.

But what is a qubit?

"Recognizing a qubit can be trickier than one might think"

Di Vincenzo, Fort. Phys. **48**, 771 (2000).

- **Qubits are subsystems** – Not necessarily the “natural” ones associated with the physical degrees of freedom – specified by the “right” **algebra of observables**, e.g.

$$\sigma_x^{(L)} = \frac{1}{6} (2 \vec{\sigma}^{(1)} \cdot \vec{\sigma}^{(2)} + \vec{\sigma}^{(2)} \cdot \vec{\sigma}^{(3)} + \vec{\sigma}^{(3)} \cdot \vec{\sigma}^{(1)}) = Ex_{1 \Leftrightarrow 2}$$

$$\sigma_y^{(L)} = -\frac{\sqrt{3}}{6} (\vec{\sigma}^{(2)} \cdot \vec{\sigma}^{(3)} - \vec{\sigma}^{(3)} \cdot \vec{\sigma}^{(1)})$$

with **both** $[\sigma_\alpha^{(L)}, \sigma_\beta^{(L)}] = 2i \epsilon^{\alpha\beta\gamma} \sigma_\gamma^{(L)}$, $\sigma_\alpha^{(L)2} = \mathbf{1}^{(L)}$

- **What else is needed ?** – Appropriate **control** for effecting both unitary and non-unitary quantum operations on subsystems. [*Dynamical control; Initialization; Read-out*].

Viola, Knill, Laflamme, quant-ph/0101090, JPA (in press).

The interaction algebra

Assume S, E initially uncorrelated. Write interaction as

$$H_{SE} = \sum_{\alpha} E_{\alpha}^{(S)} \otimes B_{\alpha}^{(E)}, \quad B_{\alpha}^{(E)} \text{ linearly independent}$$

- **Error operators** are operators that can occur in the quantum operation effected by the environment:

$$\rho \rightarrow \sum_a A_a \rho A_a^+, \quad \sum_a A_a^+ A_a = \mathbf{1}^{(S)}$$

- The possible errors are in the **interaction algebra A**:

$$A_1 = \text{span} \{ \mathbf{1}^{(S)}, E_{\alpha}^{(S)} \}$$

A = linear span of products of operators in A_1

- **Example** – Interaction algebra of **collective noise**:

$$A_1 = \text{span} \{ \mathbf{1}, J_x, J_y, J_z \}$$

A_C = algebra of **totally symmetric operators** on \mathbb{C}^8

$$\dim A_C = \dim \left(A_C^{(J=3/2)} \oplus A_C^{(J=1/2)} \right) = 4^2 + 2^2 = 20$$

Noiseless subsystems

A is an operator algebra closed under Hermitian conjugation.

- **Theorem** – The state space S of S can be represented as

$$S \simeq \bigoplus_J \mathbb{C}^{n_J} \otimes \mathbb{C}^{d_J} \quad \text{such that}$$

$$A \simeq \bigoplus_J \mathbf{1}_{n_J} \otimes \text{Mat}(d_J, \mathbb{C})$$

- Each factor \mathbb{C}^{n_J} is the state space of a n_J -dim NS of S .

A NS is defined by an **irrep** of the **commutant** A' of A :

$$A' \simeq \bigoplus_J \text{Mat}(n_J, \mathbb{C}) \otimes \mathbf{1}_{d_J}$$

- Require $n_J \geq 2$ for some J for **quantum** information encoding:
Need **non-abelian** A' . (A' contains the NS-observables!)

- **Example** – Commutant algebra of **collective noise**:

$$\begin{aligned} A'_c &= \text{algebra generated by the scalars } \{ \vec{\sigma}^{(i)} \cdot \vec{\sigma}^{(j)} \}_{i \neq j} \\ &= \text{natural representation of the group algebra } \mathbb{C} S_3 \end{aligned}$$

NSs and Error–Avoiding codes

- A non–trivial A' implies the existence of a non–trivial **group of symmetries** $G \subseteq U(A')$ such that the overall dynamics is invariant under G .
- **Example** – Collective noise:
 G = natural representation of the **permutation group** S_3 .
- A **noiseless subspace** (also: Decoherence–Free Subspace, DFS) occurs in the special case where $d_J = 1$ for some J :

$$\mathbb{C}^{n_J} \otimes \mathbb{C}^{d_J} = \mathbb{C}^{n_J} \otimes \mathbb{C} \simeq \mathbb{C}^{n_J}$$

Zanardi & Rasetti, PRL **79**, 3306 (1997).

Lidar, Chuang, Whaley, PRL **81**, 2594 (1998).

- Noiseless subspaces may be more difficult to find.
Protecting one qubit into a noiseless **subspace** against collective noise requires $n = 4$ symmetric spin $\frac{1}{2}$ particles.
NSs are the most general quantum error–avoiding codes.

NSs and Error–Suppression schemes

- Symmetries can be enforced via an **active control** action.

Viola, Knill, Lloyd, PRL **82**, 2417 (1999).

- In the limit of **fast control**, the overall **effective** Hamiltonian takes the symmetry of the control operations:

$$\bar{H}_{eff} = \frac{1}{|G|} \sum_{g_l \in G} g_l^+ H g_l \in \mathbb{C}G, \quad G = \text{decoupling group}$$

- **Example** – Permutation–invariant (collective) dynamics can be synthesized by choosing a decoupling group

$$G = \text{symmetric (permutation) group, } S_n$$

- The method allows to **dynamically generate** NSs via

$$S \simeq \oplus_J \mathbb{C}^{n_j} \otimes \mathbb{C}^{d_j}$$

$$\mathbb{C}G \simeq \oplus_J \mathbf{1}_{n_j} \otimes \text{Mat}(d_j, \mathbb{C})$$

$$\mathbb{C}G' \simeq \oplus_J \text{Mat}(n_j, \mathbb{C}) \otimes \mathbf{1}_{d_j}$$

Viola, Knill, Lloyd, PRL **85**, 3520 (2000).

QEC for general noise

- **Error analysis** – Evaluate error operators on the basis of the order in **time** with which they can occur:

$$\rho \rightarrow t \left(E_1^{(1)} \rho E_1^{(1)\dagger} + E_2^{(1)} \rho E_2^{(1)\dagger} + \dots \right) + \quad E_k^{(1)} \in A_1$$
$$t^2 \left(E_1^{(2)} \rho E_1^{(2)\dagger} + E_2^{(2)} \rho E_2^{(2)\dagger} + \dots \right) + \dots \quad E_k^{(2)} \in A_2$$

- $A_e = A_I^e = \text{span of products of } e \text{ or less operators in } A_I$
 A_e contains all error operators of **weight** (at most) e

$$A_1 \subseteq A_2 \subseteq \dots \subseteq A$$

For the independent error model, $e = \text{“number of errors”}$

- An e -Quantum Error-Correcting code is a code that can recover from **all** error operators in A_e .
- Information protected in an e -QEC code is affected by noise to order $e + 1$ or higher in time.

NSs vs QEC codes

• An *infinite*–Quantum Error–Correcting code is a code that can recover from **all** error operators in A .

• **Claim** – NS corresponds to infinite–QEC codes.

Let $S \simeq L \otimes Z \oplus D$, and $E_k(|\psi\rangle_L \otimes |s_0\rangle_Z) = |\psi\rangle_L \otimes |s_k\rangle_Z \quad \forall E_k, |s_0\rangle_Z$

$C \simeq L \otimes |s_0\rangle_Z$ satisfies the necessary & sufficient conditions for recovery from all errors in A .

Recovery super–operator $\mathbf{R} =$ “reset” operation $|s_k\rangle_Z \rightarrow |s_0\rangle_Z$
 (C, \mathbf{R}) is an infinite–QECC.

• The state $|s_0\rangle_Z$ in the “syndrome subsystem” Z is **arbitrary**:
No recovery needed for quantum **information** maintenance.
 $(L, \mathbf{R} = \mathbf{I})$ is an infinite–QECC.

• The three–spin NS is the smallest one–bit infinite–QECC for **collective errors**.

QEC codes as subsystems

- **Claim** – QEC codes can be represented as subsystems.

Use “syndrome characterization” $S \simeq C \otimes E \oplus D$.

Knill, Laflamme, PRA **55**, 900 (1997).

- **Example** – The three-bit repetition code for **bit-flip errors**:

$$(c_0|0\rangle + c_1|1\rangle) \otimes |00\rangle \rightarrow c_0|0_L\rangle + c_1|1_L\rangle = c_0|000\rangle + c_1|111\rangle$$

$$\{ E_k \} = \{ E_0 = \mathbf{1}, E_1 = X_1, E_2 = X_2, E_3 = X_3 \}$$

$$V^0 = \text{span}\{ |000\rangle, |100\rangle, |010\rangle, |001\rangle \} = \text{span}\{ |v_k^0\rangle, k=0, \dots, 3 \}$$

$$V^1 = \text{span}\{ |111\rangle, |011\rangle, |101\rangle, |110\rangle \} = \text{span}\{ |v_k^1\rangle, k=0, \dots, 3 \}$$

Then $S \simeq \mathbb{C}^8 \simeq C \otimes E$ via $|v_k^i\rangle \simeq |i\rangle_C \otimes |s_k\rangle_E$, $|s_k\rangle_E = |k_1 k_2\rangle_E$
with the identification $C \simeq C \otimes |s_0\rangle_E$ ($|s_0\rangle_E = |00\rangle_E$, no error)

Error action: $E_k \simeq \mathbf{1}^{(C)} \otimes |s_k\rangle\langle s_0|^{(E)}$

Required recovery: $R_k \simeq \mathbf{1}^{(C)} \otimes |s_0\rangle\langle s_k|^{(E)}$

QEC codes vs NSs

- C is **not** a NS of the full interaction algebra A : Repeated errors do affect the information in C , e.g. $X_1 X_2 \neq \mathbf{1}^{(C)} \otimes O(1,2)^{(E)}$.
- C **is** a NS of the algebra generated by $A_I \mathbf{R}$, where $\mathbf{R} = \{ R_a \}$ is the **reset** quantum operation on the ancillae:

$$\dots E_b, R_a, E_b R_a \simeq \mathbf{1}^{(C)} \otimes O(a,b,\dots)^{(E)}$$

- **Theorem** – Every e -error-correcting code arises as a NS of the algebra generated by $A_e \mathbf{R}$ for some quantum operation \mathbf{R} . Conversely, every NS of such an algebra $A_e \mathbf{R}$ corresponds to an e -error-correcting code.

Knill, Laflamme, Viola, PRL **84**, 2525 (2000).

- Special case $e = \infty$: Algebra generated by $A_e \mathbf{R} = A$. NSs of the full interaction algebra are infinite-QEC codes, and viceversa.

Conclusions

- The notion of a **noiseless subsystem** provides a unifying conceptual framework for realizing noise control in open quantum systems and quantum information processing.
- Experimental implementations of noiseless subsystems are within the reach of current quantum information technologies.
- Several challenges remain in theory and in practice, *e.g.*:
 - Implications of the NS notion for quantum information theory (**Operator rather than state–vector approach?**)
 - Implications of NS–encoded qubits for universal quantum computing architectures (**Encoded universality**)
 - Alternative strategy for characterizing/identifying NSs (**Generalized predictability sieve?**)
 - Improvement of attainable **quantum control** and fidelities
 - Implementation of quantum control & quantum logic over one or more NSs

..... *Work in progress*

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